### Secure Degree Equitable Dominating Graph

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#### Abstract:

Let G = (V, E) be a graph. Let  $v \in V$ : The open neighbourhood N(v) and closed neighbourhood N[v] are defined by N(v) = $\{u \in V : uv \in E\}$  and  $N[v] = N(v) \cup \{v\}$ 

A set contains D is the secure degree equitable dominating set in G it satisfies the following conditions

- i) A vertex  $u \in V$  is set to be degree equitable with a vertex  $v \in V$  if  $|deg(u) - deg(v)| \le 1$ .
- ii) A dominating set S of G is a secure dominating set if for each  $u \in V$  - S there exists  $v \in N(u) \cap S$  such that  $(S - \{v\}) \cup \{u\}$  is a secure dominating set.

The minimum cardinality of the secure degree equitable dominating set  $D_{sde}(G)$  and is denoted by  $\delta_s^{de}(G)$ . This paper find the secure degree equitable domination number of  $\delta_s^{de}(G)$  of cycle graphs, path graphs and complete graphs and also find the

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### I. Introduction:

Let G = (V, E) be a graph, mean a finite undirected graph with neither loops nor multiple edges.

Let G = (V, E) be a graph, it is said to be complete all the vertices of V are adjacent to each other.

Degree of a vertex  $v \in V$  can be defined number of edges incident with v, it can be denoted by deg(v).

Let G = (V, E) be a graph. Let  $v \in V$ : The open neighbourhoodN(v) and closed neighbourhood N[v] are defined by N(v) =  $\{u \in V : uv \in E\}$  and N[v] = N(v) U  $\{v\}$ 

#### **II. Theorem and Proof:**

The set S is called a dominating set of G if every vertex in V - S is adjacent to at least one vertex in S.

A dominating set S of

G is a secure dominating set if for each  $u \in V$  - S there exists

 $v \in N(u) \cap S$  such that  $(S - \{v\}) \cup \{u\}$  is a dominating set.

A vertex  $u \in V$  is set to be degree equitable with a vertex  $v \in V$  if

 $|\deg(u) - \deg(v)| \le 1.$ 

A subset *D* of *V* is called an equitable dominating set if for every  $v \in V - D$  there exists *a* vertex  $u \in D$  such that  $uv \in E(G)$  and  $|\deg(u) - \deg(v)| \le 1$ .

A set contains D is the secure degree equitable dominating set in G it satisfies the following conditions

- iii) A vertex  $u \in V$  is set to be degree equitable with a vertex  $v \in V$  if  $|deg(u) - deg(v)| \le 1$ .
- iv) A dominating set S of G is a secure dominating set if for each  $u \in V$  - S there exists  $v \in N(u) \cap S$  such that  $(S - \{v\}) \cup \{u\}$  is a secure dominating set.

The minimum cardinality of the secure degree equitable dominating set is called secure degree equitable domination number and is denoted by  $\delta_s^{de}$ .

*Theorem 1:* Let  $K_n$  be a complete graph of order  $m \ge 2$ , then  $\delta_s^{de}(K_m) = 1$ .

Proof: Given a secure degree equitable doming set  $D_{sde}$  of  $K_m = \{x_1, x_2, x_3, \dots, x_m\}$ , assume that  $S = \{x_1\} \in D_{sde}$ , Then  $x_1$  dominates all other vertices in  $K_m$ , and since  $K_m$  is a complete graph, then for every  $x_{i \neq 1} \in K_m$ ,  $x_1$  and  $x_{i \neq 1}$ are adjacent and

 $|\deg(\mathbf{x}_1) - \deg(\mathbf{x}_{i \neq 1})| = 0$ Hence

 $|\deg(\mathbf{x}_1) - \deg(\mathbf{x}_{i \neq 1})| \le 1$ 

For each

 $x_{i \neq 1} \in V - S = \{x_2, x_3, \dots, x_m\}$  there exists  $x_{i \neq 1} \in N(x_1) \cap S$  such that

 $(S - \{x_{i \neq 1}\}) \cup \{x_1\}$  is a secure dominating set.

Then  $D_{sde}(K_m) = \{x_1\}$  which gives  $\delta_s^{de}(K_m) = 1$ .

For example

 $K_4$  is a complete graph (in Figure 1) with V = {a, b, c, d}.

Let  $S = \{a\}$  and V-  $S = \{b, c, d\}$  and

 $| \ deg \ (a) \ - \ \{deg(b)or \ deg(c) \ or \ deg(d)\}|{=}0.$  Therefore  $K_4$  is degree equitable.

For each x (say d) belongs to V- S = {b, c, d} and N (d) = {a, c, b} and N (b)  $\cap$  S = {d} and then (S-{a})  $\cup$  {b} = {b, c, d} is a secure dominating set.  $\delta_s^{de}$  (K<sub>4</sub>) =1.



Figure 1

Proof:

For n=2,

degree equitable.

K<sub>4</sub> (Removing a vertex {a} from S and inserting the adjacent vertex d to S. Therefore K<sub>4</sub> is secure degree equitable domination. $\delta_s^{de}$  (K<sub>4</sub>) =1.)

Theorem 2: Let  $P_n$  be a path graph of order  $n \ (n \ge 2)$ , then  $\delta_s^{de}$  ( $P_n$ ) is greatest integer function of  $(\frac{n}{2})$ .



a





Figure 2 In P<sub>2</sub> (In Figure 2),  $S = \{a\}$ .Removing a vertex  $\{a\}$  from S and inserting the adjacent vertex  $\{b\}$ to S. Then  $S = \{b\}$  is a secure dominating

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b

P<sub>3</sub>

set. Therefore  $\delta_s^{de}$  (P<sub>2</sub>) is greatest integer function of  $\left(\frac{2}{2}\right) = 1$ .

Every vertex in  $P_n$  degree is either 2(between

vertices) or 1(endvertex). Therefore clearly  $P_n$  is





Figure 3 In  $P_3($ in Figure 3),  $S = \{a, b\}$ . Removing a vertex  $\{b\}$  from S and inserting the adjacent

For n = 4,

vertex {c} to S. Then S={c, b} is a secure dominating set. Therefore  $\delta_s^{de}$  (P<sub>3</sub>) is greatest integer function of  $\left(\frac{3}{2}\right) = 2$ .





Figure 4

In  $P_4$  (Figure 4),  $S = \{a, c\}$ . Removing a vertex  $\{a\}$  from S and inserting the adjacent vertex  $\{b\}$  to S. Then  $S=\{c, b\}$  is a secure dominating set.

Therefore  $\delta_s^{de}$  (P<sub>4</sub>) is greatest integer function of  $\left(\frac{4}{2}\right) = 2$ .



Figure 5

In  $P_5$  (in Figure 5),  $S = \{a, c, f\}$ . Removing a vertex  $\{f\}$  from S and inserting the adjacent vertex  $\{d\}$  to S. Then  $S=\{a, c, d\}$  is a secure

dominating set. Therefore  $\delta_s^{de}$  (P<sub>5</sub>) is greatest integer function of  $\left(\frac{5}{2}\right) = 3$ .

For n = 6,



Figure 6

In P<sub>6</sub>, (in Figure 6) S = {a, c, f}. Removing a vertex {f} from S and inserting the adjacent vertex {g} to S. Then S={a, c, g} is a secure dominating set.Therefore  $\delta_s^{de}$  (P<sub>6</sub>) is greatest integer function of  $\left(\frac{6}{2}\right) = 3$ .

Therefore in general  $P_n$  be a path graph of order n, then  $\delta_s^{de}$  (P<sub>n</sub>) is greatest integer function of  $(\frac{n}{2})$ .

### Proof:

Theorem 3: Let  $C_n$  be a complete graph of order  $n \ge 3$ , then  $\delta_s^{de}(C_n)$  is greatest integer function  $of(\frac{n}{2})$ .

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Every vertex in  $C_n$  is degree of 2. Therefore clearly  $C_n$  is degree equitable. For n = 3,





In C<sub>3</sub> (In Figure 7) , S = {a, c}. Removing a vertex {a} from S and inserting the adjacent vertex {b} to S. Then

S={c, b} is a secure dominating set. Therefore  $\delta_s^{de}$  (C<sub>3</sub>) is greatest integer function of  $\left(\frac{3}{2}\right) = 2$ . For n = 4,



Figure 8

In C<sub>4</sub> (In Figure 8) , S =  $\{a, c\}$ . Removing a vertex  $\{a\}$  from S and inserting the adjacent vertex  $\{b\}$  to S. Then

S={c, b} is a secure dominating set. Therefore  $\delta_s^{de}$  (C<sub>4</sub>) is greatest integer function of  $\left(\frac{4}{2}\right) = 2$ . For n = 5,



Figure 9

In P<sub>5</sub> (In Figure 9)

, S = {a, c, e}. Removing a vertex {e} from S and inserting the adjacent vertex {d} to S. Then S={a, c, d} is a secure dominating set. Therefore  $\delta_s^{de}$  (P<sub>5</sub>) is greatest integer function of  $\left(\frac{5}{2}\right) = 3$ .

For n = 6,



Figure 10

In C<sub>6</sub> (In Figure 10)

, S = {a, c, f}. Removing a vertex {f} from S and inserting the adjacent vertex {e} to S. Then S={a, c, e} is a secure dominating set. Therefore  $\delta_s^{de}$  (C<sub>6</sub>) is greatest integer function of  $\left(\frac{6}{2}\right) = 3$ .

#### **III. Conclusion:**

From this paper we can able to find the  $\delta_s^{de}(G)$  of cycle graphs, path graphs and complete graphs and also find the secure degree equitable domination number.

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